

# Properties of the intuitionistic fuzzy implication $\rightarrow_{186}$

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**Abstract:** In [8], a new Fodor’s type of intuitionistic fuzzy implication, numbered as  $\rightarrow_{186}$ , was defined and some of its properties were studied. The present paper is a continuation of the previous one. New interesting properties of implication  $\rightarrow_{186}$  are formulated and checked.

**Keywords:** Implication, Intuitionistic fuzzy implication, Intuitionistic fuzzy logic.

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## 1 Introduction

In [7, 15, 16] some Fodor’s type of intuitionistic fuzzy implications were introduced. The first five were included in [6], while the latest one, numbered as  $\rightarrow_{186}$ , was introduced latter – in [8]. There, some of its basic properties were studied, e.g., it was checked which axioms of Klir and Yuan, of Kolmogorov, of Łukasiewicz–Tarski and of intuitionistic logic axioms are valid. Here, we continue this research, studying new properties of intuitionistic fuzzy implication  $\rightarrow_{186}$ .

In the intuitionistic fuzzy logic (see [1, 2, 4, 5], each proposition, variable or formula is evaluated with two degrees – “truth degree” or “degree of validity” and “falsity degree” or “degree

of non-validity". Thus, to each one of these objects, e.g.,  $p$ , two real numbers,  $\mu(p)$  and  $\nu(p)$ , are assigned with the following constraint:

$$\mu(p), \nu(p) \in [0, 1] \text{ and } \mu(p) + \nu(p) \leq 1.$$

Let

$$\pi(p) = 1 - \mu(p) - \nu(p).$$

This function determines the degree of uncertainty (indeterminacy).

Let an evaluation function  $V$  be defined over a set of propositions  $\mathcal{S}$ , in such a way that for  $p \in \mathcal{S}$ :

$$V(p) = \langle \mu(p), \nu(p) \rangle.$$

Hence the function  $V : \mathcal{S} \rightarrow [0, 1] \times [0, 1]$  gives the truth and falsity degrees of all elements of  $\mathcal{S}$  – the set of logical objects that we use (in general case – formulas).

We assume that the evaluation function  $V$  assigns to the logical truth  $T$ :

$$V(T) = \langle 1, 0 \rangle,$$

and to the logical falsity it assigns the value:  $F$

$$V(F) = \langle 0, 1 \rangle.$$

As it was discussed [3], the first two intuitionistic fuzzy negations are

$$V(\neg_1 p) = \langle \nu(p), \mu(p) \rangle,$$

$$V(\neg_2 p) = \langle \overline{\text{sg}}(\mu(p)), \text{sg}(\mu(p)) \rangle,$$

where here and below

$$\text{sg}(x) = \begin{cases} 1 & \text{if } x > 0 \\ 0 & \text{if } x \leq 0 \end{cases},$$

and

$$\overline{\text{sg}}(x) = \begin{cases} 0 & \text{if } x > 0 \\ 1 & \text{if } x \leq 0 \end{cases}.$$

The second negation does not satisfy the equality  $V(\neg_2 \neg_2 p) = V(p)$ , and, as we see below, it exhibits truly intuitionistic behaviour. Here, we define only the operations “disjunction”, “conjunction” and “implication”, originally introduced in [1], that have classical logic analogues, as follows:

$$V(p \vee q) = \langle \max(\mu(p), \mu(q)), \min(\nu(p), \nu(q)) \rangle,$$

$$V(p \wedge q) = \langle \min(\mu(p), \mu(q)), \max(\nu(p), \nu(q)) \rangle,$$

$$V(p \rightarrow_{186} q) = \langle \max(\nu(p), \mu(q)), \min(\nu(p), \mu(q)) \rangle.$$

For the needs of the discussion below, we define the notions of Intuitionistic Fuzzy Tautology (IFT, see, e.g. [1, 3]) and tautology.

Formula  $A$  is an IFT if and only if (iff) for every evaluation function  $V$ , if  $V(A) = \langle a, b \rangle$ , then,

$$a \geq b,$$

while it is a (classical) tautology if and only if for every evaluation function  $V$ , if  $V(A) = \langle a, b \rangle$ , then,

$$a = 1, b = 0.$$

Below, when it is clear, we will omit notation “ $V(A)$ ”, using directly “ $A$ ” of the intuitionistic fuzzy evaluation of  $A$ . Also, for brevity, in a lot of places, instead of the IFP  $\langle \mu(A), \nu(A) \rangle$  we will use the IFP  $\langle a, b \rangle$ , where  $a, b \in [0, 1]$  and  $a + b \leq 1$ .

It is also suitable, if  $\langle a, b \rangle$  and  $\langle c, d \rangle$  are IFPs, to have

$$\langle a, b \rangle \leq \langle c, d \rangle \text{ iff } a \leq c \text{ and } b \geq d$$

and

$$\langle a, b \rangle \geq \langle c, d \rangle \text{ iff } a \geq c \text{ and } b \leq d.$$

## 2 Properties of the intuitionistic fuzzy implication $\rightarrow_{186}$

In [8], the following intuitionistic fuzzy implication was introduced:

$$V(x \rightarrow_{186} y) = \langle \overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), \text{sg}(d - b) \min(a, d) \rangle$$

and it is shown that it generates the standard (classical) intuitionistic fuzzy negation  $\neg_1$ , or briefly  $\neg$ .

G. F. Rose’s formula [13, 14] has the form:

$$((\neg\neg A \rightarrow_{186} A) \rightarrow_{186} (\neg\neg A \vee \neg A)) \rightarrow_{186} (\neg\neg A \vee \neg A). \quad (1)$$

For it, the following two theorems are valid.

**Theorem 1.** Implication  $\rightarrow_{186}$  satisfies Rose’s formula (1) as a tautology.

*Proof.* Let  $A$  be a formula for which  $V(A) = \langle a, b \rangle$ . Then

$$\begin{aligned} & V(((\neg\neg A \rightarrow_{186} A) \rightarrow_{186} (\neg\neg A \vee \neg A)) \rightarrow_{186} (\neg\neg A \vee \neg A)) \\ &= ((\neg\neg \langle a, b \rangle \rightarrow_{186} \langle a, b \rangle) \rightarrow_{186} (\neg\neg \langle a, b \rangle \vee \neg \langle a, b \rangle)) \rightarrow_{186} (\neg\neg \langle a, b \rangle \vee \neg \langle a, b \rangle) \\ &= ((\langle a, b \rangle \rightarrow_{186} \langle a, b \rangle) \rightarrow_{186} (\langle a, b \rangle \vee \langle b, a \rangle)) \rightarrow_{186} (\langle a, b \rangle \vee \langle b, a \rangle) \\ &= (\langle \overline{\text{sg}}(b - b) + \text{sg}(b - b) \max(b, a), \text{sg}(b - b) \min(a, b) \rangle \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle) \\ &\quad \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle \\ &= (\langle 1, 0 \rangle \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle) \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle \\ &= \langle \overline{\text{sg}}(\min(a, b) - 0) + \text{sg}(\min(a, b) - 0) \max(0, \max(a, b)), \text{sg}(\min(a, b) - 0) \min(1, \min(a, b)) \rangle \\ &\quad \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle \end{aligned}$$

$$\begin{aligned}
&= \langle \overline{\text{sg}}(\min(a, b)) + \text{sg}(\min(a, b)) \max(a, b), \text{sg}(\min(a, b)) \min(a, b) \rangle \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle \\
&\text{(because for each real number } x \geq 0: \text{sg}(x)x = x.) \\
&= \langle \overline{\text{sg}}(\min(a, b)) + \text{sg}(\min(a, b)) \max(a, b), \min(a, b) \rangle \rightarrow_{186} \langle \max(a, b), \min(a, b) \rangle \\
&= \langle \overline{\text{sg}}(\min(a, b) - \min(a, b)) + \text{sg}(\min(a, b) - \min(a, b)) \max(\min(a, b), \max(a, b)), \\
&\quad \text{sg}(\min(a, b) - \min(a, b)) \min(\overline{\text{sg}}(\min(a, b)) + \text{sg}(\min(a, b)) \max(a, b), \min(a, b)) \rangle \\
&= \langle \overline{\text{sg}}(0) + \text{sg}(0) \max(\min(a, b), \max(a, b)), \\
&\quad \text{sg}(0) \min(\overline{\text{sg}}(\min(a, b)) + \text{sg}(\min(a, b)) \max(a, b), \min(a, b)) \rangle \\
&= \langle 1, 0 \rangle.
\end{aligned}$$

Therefore, Rose's formula (1) is a tautology.  $\square$

**Corollary 1.** Implication  $\rightarrow_{186}$  satisfies Rose's formula (1) as an intuitionistic fuzzy tautology.

The next assertions are proved in a similar manner and by this reason, we omit their proofs. Now, we discuss the following formulas, inspired by (1):

$$(A \vee \neg A) \rightarrow_{186} (A \rightarrow_{186} \neg \neg A) \quad (2)$$

$$(\neg \neg A \vee \neg A) \rightarrow_{186} (A \rightarrow_{186} \neg \neg A) \quad (3)$$

$$(A \rightarrow_{186} \neg \neg A) \rightarrow_{186} (A \vee \neg A) \quad (4)$$

$$(A \rightarrow_{186} \neg \neg A) \rightarrow_{186} (\neg \neg A \vee \neg A) \quad (5)$$

Obviously, in the classical propositional calculus, all these four formulas are tautologies. Now, we study their properties in the intuitionistic fuzzy case.

**Theorem 2.** Implication  $\rightarrow_{186}$  satisfies (2) and (3) as tautologies.

**Corollary 2.** Implication  $\rightarrow_{186}$  satisfies (2) and (3) as intuitionistic fuzzy tautologies.

**Theorem 3.** Implication  $\rightarrow_{186}$  satisfies (4) and (5) as intuitionistic fuzzy tautologies, but not as tautologies.

Now, following [4], we discuss the well-known Contraposition Law

$$(A \rightarrow_{186} B) \rightarrow_{186} (\neg B \rightarrow_{186} \neg A) \quad (6)$$

and its modified version

$$(\neg \neg A \rightarrow_{186} \neg \neg B) \rightarrow_{186} (\neg B \rightarrow_{186} \neg A). \quad (7)$$

For them, the following assertions are valid.

**Theorem 4.** Implication  $\rightarrow_{186}$  satisfies (6) and (7) as intuitionistic fuzzy tautologies, but not as tautologies.

*Proof.* Let  $A$  and  $B$  be a formulas for which  $V(A) = \langle a, b \rangle$  and  $V(B) = \langle c, d \rangle$ . Then for (6) we obtain

$$\begin{aligned}
& V((A \rightarrow_{186} B) \rightarrow_{186} (\neg B \rightarrow_{186} \neg A)) \\
&= (\langle a, b \rangle \rightarrow_{186} \langle c, d \rangle) \rightarrow_{186} (\neg \langle c, d \rangle \rightarrow_{186} \neg \langle a, b \rangle) \\
&= (\langle a, b \rangle \rightarrow_{186} \langle c, d \rangle) \rightarrow_{186} (\langle d, c \rangle \rightarrow_{186} \langle b, a \rangle) \\
&= \langle \overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), \text{sg}(d - b) \min(a, d) \rangle \\
&\rightarrow_{186} \langle \overline{\text{sg}}(a - c) + \text{sg}(a - c) \max(c, b), \text{sg}(a - c) \min(a, d) \rangle \\
&= \langle \overline{\text{sg}}(\text{sg}(a - c) \min(a, d) - \text{sg}(d - b) \min(a, d)) + \text{sg}(\text{sg}(a - c) \min(a, d) \\
&\quad - \text{sg}(d - b) \min(a, d)) \cdot \max(\text{sg}(d - b) \min(a, d), \overline{\text{sg}}(a - c) + \text{sg}(a - c) \max(c, b)), \\
&\quad \text{sg}(\text{sg}(a - c) \min(a, d) - \text{sg}(d - b) \min(a, d)) \min(\overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), \\
&\quad \text{sg}(a - c) \min(a, d)) \rangle.
\end{aligned}$$

Let

$$\begin{aligned}
X &\equiv \overline{\text{sg}}(\text{sg}(a - c) \min(a, d) - \text{sg}(d - b) \min(a, d)) + \text{sg}(\text{sg}(a - c) \min(a, d) - \text{sg}(d - b) \min(a, d)) \\
&\quad \cdot \max(\text{sg}(d - b) \min(a, d), \overline{\text{sg}}(a - c) + \text{sg}(a - c) \max(c, b)) - \text{sg}(\text{sg}(a - c) \min(a, d) \\
&\quad - \text{sg}(d - b) \min(a, d)) \cdot \min(\overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), \text{sg}(a - c) \min(a, d)).
\end{aligned}$$

1. If  $a > c$ , then

$$\begin{aligned}
X &= \overline{\text{sg}}(\min(a, d) - \text{sg}(d - b) \min(a, d)) + \text{sg}(\min(a, d) - \text{sg}(d - b) \min(a, d)) \\
&\quad \cdot \max(\text{sg}(d - b) \min(a, d), \max(c, b)) - \text{sg}(\min(a, d) - \text{sg}(d - b) \min(a, d)) \\
&\quad \cdot \min(\overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), \min(a, d)).
\end{aligned}$$

1.1. If  $d > b$ , then

$$\begin{aligned}
X &= \overline{\text{sg}}(\min(a, d) - \min(a, d)) + \text{sg}(\min(a, d) - \min(a, d)) \cdot \max(\min(a, d), \max(c, b)) \\
&\quad - \text{sg}(\min(a, d) - \min(a, d)) \cdot \min(\max(b, c), \min(a, d)) \\
&= 1 - 0 = 1 > 0.
\end{aligned}$$

1.2. If  $d \leq b$ , then

$$\begin{aligned}
X &= \overline{\text{sg}}(\min(a, d)) + \text{sg}(\min(a, d)) \cdot \max(c, b) - \text{sg}(\min(a, d)) \cdot \min(1, \min(a, d)) \\
&= \overline{\text{sg}}(\min(a, d)) + \text{sg}(\min(a, d)) \cdot \max(c, b) - \text{sg}(\min(a, d)) \cdot \min(a, d).
\end{aligned}$$

1.2.1. From  $a > c$  it follows that  $a > 0$ . Let  $d = 0$ . Then

$$X = \overline{\text{sg}}(0) + \text{sg}(0) \cdot \max(c, b) - \text{sg}(0) \cdot \min(a, d) = 1 > 0.$$

1.2.2. Let  $d > 0$ . Then

$$X = 0 + \max(c, b) - \min(a, d) \geq b - d \geq 0.$$

2. If  $a \leq c$ , then

$$X = \overline{\text{sg}}(-\text{sg}(d - b) \min(a, d)) + \text{sg}(-\text{sg}(d - b) \min(a, d))$$

$$\cdot \max(\text{sg}(d - b) \min(a, d), 1) - \text{sg}(0 - \text{sg}(d - b) \min(a, d)) \cdot \min(\overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), 0).$$

$$X = 1 - 0 \cdot \min(\overline{\text{sg}}(d - b) + \text{sg}(d - b) \max(b, c), 0) = 1 > 0.$$

Therefore, in all cases  $X \geq 0$ , i.e., (6) is an IFT, but in case 1.2.2, the expression is not 1, as in the rest cases, i.e., (6) is not a tautology. The validity of (7) follows directly from (6).  $\square$

### 3 Conclusion

In next research other properties of the implications will be introduced and studied. All they show that intuitionistic fuzzy sets and logics in the sense, described in [2, 3], correspond to the ideas of Brouwer's intuitionism (see [9, 11, 12]).

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