

INTUITIONISTIC FUZZY INTERPRETATION OF KUN'S AXIOM

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In [1] V. Kostjuk formalize the Kun's axiom "if  $A$  and  $B$  are two theories from which follow contradictory consequences, then the expression containing the conjunction of  $A$  and  $B$  is false" in the form

$$((A \supset C) \& (B \supset \neg C)) \supset \neg(A \& B).$$

He assert that this axiom will have more correct form, if it contains modalities and he offers the axiom:

$$(\Box(A \supset C) \& \Box(B \supset \neg C)) \supset \neg\Diamond(A \& B).$$

We shall show that these axioma can be interpreted as Intuitionistic Fuzzy Tautologies (IFTs).

First, following [2,3] we shall give some short remarks on the intuitionistic fuzzy logic related to the below text.

To each proposition (see, e.g., [4]) we can assign its truth value: truth – denoted by 1, or falsity – 0. In the case of fuzzy logic this truth value is a real number in the interval  $[0, 1]$  and may be called "truth degree" of a particular proposition. Here we add one more value – "falsity degree" – which will be in the interval  $[0, 1]$  as well. Thus two real numbers,  $\mu(p)$  and  $\nu(p)$ , are assigned to the proposition  $p$  with the following constraint to hold:

$$\mu(p) + \nu(p) \leq 1.$$

Let this assignment be provided by an evaluation function  $V$  defined over a set of propositions  $S$  in such a way that:

$$V(p) = \langle \mu(p), \nu(p) \rangle.$$

Hence the function  $V : S \rightarrow [0, 1] \times [0, 1]$  gives the truth and falsity degrees of all propositions in  $S$ .

When the values  $V(p)$  and  $V(q)$  of the propositions  $p$  and  $q$  are known, the evaluation function  $V$  can be extended also for the operations “&”, “ $\vee$ ” through the definition :

$$\begin{aligned} V(\neg p) &= \langle \nu(p), \mu(p) \rangle, \\ V(p \& q) &= \langle \min(\mu(p), \mu(q)), \max(\nu(p), \nu(q)) \rangle, \\ V(p \supset q) &= \langle \max(\nu(p), \mu(q)), \min(\mu(p), \nu(q)) \rangle, \end{aligned}$$

As with the operations over IFS, it will be convenient to define for the propositions  $p, q \in S$ :

$$\begin{aligned} \neg V(p) &= V(\neg p), \\ V(p) \wedge V(q) &= V(p \& q), \\ V(p) \rightarrow V(q) &= V(p \supset q). \end{aligned}$$

For the needs of the discussion below we shall define the notion of IFT through:

$$\text{“}A \text{ is an IFT” iff if } V(A) = \langle a, b \rangle, \text{ then } a \geq b.$$

Now, we shall prove the following

**THEOREM:** For every three propositional forms  $A, B$  and  $C$ :

- (a)  $((A \supset C) \& (B \supset \neg C)) \supset \neg(A \& B)$ ,
  - (b)  $(\square(A \supset C) \& \square(B \supset \neg C)) \supset \neg \diamond(A \& B)$
- are IFTs.

Proof: Let  $V(A) = \langle \mu_A, \nu_A \rangle$ ,  $V(B) = \langle \mu_B, \nu_B \rangle$ ,  $V(C) = \langle \mu_C, \nu_C \rangle$

$$\begin{aligned} \text{(a)} \quad & V(((A \supset C) \& (B \supset \neg C)) \supset \neg(A \& B)) \\ &= ((\langle a, b \rangle \rightarrow \langle e, f \rangle) \wedge (\langle c, d \rangle \rightarrow \langle f, e \rangle)) \rightarrow \neg \langle a, b \rangle \wedge \langle c, d \rangle \\ &= (\langle \max(b, e), \min(a, f) \rangle \wedge \langle \max(d, f), \min(c, e) \rangle) \\ &\quad \rightarrow \neg \langle \min(a, c), \max(b, d) \rangle \\ &= \langle \min(\max(b, e), \max(d, f)), \max(\min(a, f), \min(c, e)) \rangle \\ &\quad \rightarrow \langle \max(b, d), \min(a, c) \rangle \\ &= \langle \max(b, d, \min(a, f), \min(c, e)), \min(a, c, \max(b, e), \max(d, f)) \rangle. \end{aligned}$$

Let

$$X = \max(b, d, \min(a, f), \min(c, e)) - \min(a, c, \max(b, e), \max(d, f)).$$

If  $a \leq f$ , then

$$X \geq \max(b, d, a) - a \geq a - a = 0$$

If  $a > f$ , then

$$X \geq \max(b, d, f) - \max(d, f) \geq 0.$$

Therefore, the axiom is an IFT.

- (b)  $V(\square(A \supset C) \& \square(B \supset \neg C)) \supset \neg \diamond(A \& B)$

$$\begin{aligned}
&= (\Box \langle \max(b, e), \min(a, f) \rangle \wedge \Box \langle \max(d, f), \min(c, e) \rangle) \\
&\quad \rightarrow \neg \Diamond \langle \min(a, c), \max(b, d) \rangle \\
&= (\langle \max(b, e), 1 - \max(b, e) \rangle \wedge \langle \max(d, f), 1 - \max(d, f) \rangle) \\
&\quad \rightarrow \neg \langle 1 - \max(b, d), \max(b, d) \rangle \\
&= \langle \min(\max(b, e), \max(d, f)), \max(1 - \max(b, e), 1 - \max(d, f)) \rangle \\
&\quad \rightarrow \langle \max(b, d), 1 - \max(b, d) \rangle \\
&= \langle \max(b, d, 1 - \max(b, e), 1 - \max(d, f)), \min(1 - \max(b, d), \max(b, e), \max(d, f)) \rangle.
\end{aligned}$$

Let

$$X = \max(b, d, 1 - \max(b, e), 1 - \max(d, f)) - \min(1 - \max(b, d), \max(b, e), \max(d, f)).$$

If  $d \geq e$ , then

$$X \geq \max(b, d) - \max(b, e) \geq 0$$

Let  $d < e$  and let  $b \geq e$ , then

$$X \geq b - \max(b, e) = b - b = 0.$$

Let  $d < e$  and  $b < e$ , then

$$X = \max(b, d, 1 - e, 1 - \max(d, f)) - \min(1 - \max(b, d), e, \max(d, f)).$$

If  $d \geq f$ , then

$$X \geq 1 - d - 1 + \max(b, d) = \max(b, d) - d \geq 0.$$

If  $d < f$ , then

$$X = 1 - f - e \geq 0.$$

Therefore, the axiom is an IFT.

#### References:

- [1] Kostjuk V., Elements of Modal Logic, Kiev, Naukova Dumka, 1978 (in Russian).
- [2] Atanassov K. Two variants of intuitionistic fuzzy propositional calculus. Preprint IM-MFAIS-5-88, Sofia, 1988.
- [3] Atanassov K., Gargov G. C. R. Acad. bulg. sci., Tome 43, 1990, No 3, 9-12.
- [4] Mendelson E., Introduction to Mathematical Logic, Princeton, NJ, 1964.