

Intuitionistic fuzzy implication \rightarrow_{189}

Lilija Atanassova

Institute of Information and Communication Technologies
Bulgarian Academy of Sciences
Acad. G. Bonchev Str., Bl. 2, Sofia 1113, Bulgaria
e-mail: l.c.atanassova@gmail.com

Abstract: In [4], some new intuitionistic fuzzy operations are defined and their properties are studied. On the basis of the third of them, a new intuitionistic fuzzy implication is introduced here, numbered as \rightarrow_{189} and some of its properties will be studied.

Keywords: Implication, Intuitionistic fuzzy implication, Intuitionistic fuzzy logic.

AMS Classification: 03E72.

1 Introduction

In the present paper, we continue our research, related to some new intuitionistic fuzzy implications. Now, we use results from [4], where five new intuitionistic fuzzy operations, including multiplication, were introduced.

As it is discussed in [1], each proposition, variable or formula is evaluated with two degrees – “truth degree” or “degree of validity”, and “falsity degree” or “degree of non-validity”. Thus, to each of these objects, e.g., p , two real numbers, $\mu(p)$ and $\nu(p)$, are assigned with the following constraint:

$$\mu(p), \nu(p) \in [0, 1] \text{ and } \mu(p) + \nu(p) \leq 1.$$

In [2], the object $\langle \mu(p), \nu(p) \rangle$ is introduced under the name Intuitionistic Fuzzy Pair (IFP). Formula A is an Intuitionistic Fuzzy Tautology (IFT) if and only if (*iff*) for every evaluation function V , if $V(A) = \langle a, b \rangle$, then, $a \geq b$, while it is a (classical) tautology *iff* for every evaluation function V , if $V(A) = \langle a, b \rangle$, then, $a = 1, b = 0$.

Following [2], we will mention that, if an IFP is an IFT, we call it Intuitionistic Fuzzy Tautological Pair (IFTP) and if it is a tautology – Tautological Pair (TP).

In [1], different intuitionistic fuzzy operations are introduced, e.g., intuitionistic fuzzy disjunction, conjunction and (classical) negation, by

$$V(p \vee q) = \langle \max(\mu(p), \mu(q)), \min(\nu(p), \nu(q)) \rangle,$$

$$V(p \wedge q) = \langle \min(\mu(p), \mu(q)), \max(\nu(p), \nu(q)) \rangle,$$

$$V(\neg p) = \langle \nu(p), \mu(p) \rangle.$$

Below, when it is clear, we will omit notation “ $V(A)$ ”, using directly “ A ” of the intuitionistic fuzzy evaluation of A . Also, for brevity, in a lot of places, instead of the IFP $\langle \mu(A), \nu(A) \rangle$ we use the IFP $\langle a, b \rangle$, where $a, b \in [0, 1]$ and $a + b \leq 1$.

In [4], for two IFPs $x = \langle a, b \rangle$ and $y = \langle c, d \rangle$ are introduced five novel operations from multiplicative type. Here, we use only one of them:

$$x \times_3 y = \langle ac, bd \rangle$$

and on its basis we introduce a new intuitionistic fuzzy implication. In some sense, it is analogous of implication $\rightarrow_{\text{@}}$, that in [1] was denoted by \rightarrow_{139} .

For operation \times_3 , in [4] it was checked that it was defined correctly and for the above x and y and for $z = \langle e, f \rangle$:

$$x \times_3 y = y \times_3 x,$$

$$(x \times_3 y) \times_3 z = x \times_3 (y \times_3 z).$$

We can see that

$$\langle 0, 1 \rangle \times_3 \langle 0, 1 \rangle = \langle 0, 1 \rangle,$$

$$\langle 0, 1 \rangle \times_3 \langle 0, 0 \rangle = \langle 0, 0 \rangle,$$

$$\langle 0, 1 \rangle \times_3 \langle 1, 0 \rangle = \langle 0, 0 \rangle,$$

$$\langle 0, 0 \rangle \times_3 \langle 0, 1 \rangle = \langle 0, 0 \rangle,$$

$$\langle 0, 0 \rangle \times_3 \langle 0, 0 \rangle = \langle 0, 0 \rangle,$$

$$\langle 0, 0 \rangle \times_3 \langle 1, 0 \rangle = \langle 0, 0 \rangle,$$

$$\langle 1, 0 \rangle \times_3 \langle 1, 0 \rangle = \langle 1, 0 \rangle,$$

$$\langle 1, 0 \rangle \times_3 \langle 0, 0 \rangle = \langle 0, 0 \rangle,$$

$$\langle 1, 0 \rangle \times_3 \langle 0, 1 \rangle = \langle 0, 0 \rangle,$$

and for each IFP x :

$$\langle 0, 1 \rangle \times_3 x = \langle 0, b \rangle = x \times_3 \langle 0, 1 \rangle,$$

$$\langle 0, 0 \rangle \times_3 x = \langle 0, 0 \rangle = x \times_3 \langle 0, 0 \rangle,$$

$$\langle 1, 0 \rangle \times_3 x = \langle a, 0 \rangle = x \times_3 \langle 1, 0 \rangle.$$

In [4] it is proved that if

$$\mathcal{L} = \{\langle a, b \rangle \mid a, b \in [0, 1] \ \& \ a + b \leq 1\}$$

is the set of all IFPs, then $\langle \mathcal{L}, \times_3 \rangle$ is a commutative semi-group and if x and y are IFTPs, then $x \times_3 y$, is an IFTP, while if x and y are TPs, then $x \times_3 y$ is a TP.

Using the classical negation, defined in intuitionistic fuzzy propositional logic, for \times_3 we obtain for every two IFPs x and y :

$$\neg(\neg x \times_3 \neg y) = \neg(\neg \langle a, b \rangle \times_3 \neg \langle c, d \rangle)$$

$$\neg(\langle b, a \rangle \times_3 \langle d, c \rangle) = \neg \langle bd, ac \rangle$$

$$\langle ac, bd \rangle = x \times_3 y.$$

Therefore, operation \times_3 , similarly to operation \textcircled{a} (see [1]), simultaneously has the behaviour of both the operations disjunction and conjunction.

2 Intuitionistic fuzzy implication \rightarrow_{189} and its properties

In the paper, using the standard logical formula

$$x \rightarrow y = \neg x \vee y,$$

we obtain the new intuitionistic fuzzy implication

$$x \rightarrow_{189} y = \neg x \vee y = \langle bc, ad \rangle.$$

First, we see that

$$0 \leq bc + ad \leq b + a \leq 1,$$

i.e., implication \rightarrow_{189} is defined correctly.

Second, we see that

$$\langle 0, 1 \rangle \rightarrow_{189} \langle 0, 1 \rangle = \langle 0, 0 \rangle,$$

$$\langle 0, 1 \rangle \rightarrow_{189} \langle 1, 0 \rangle = \langle 1, 0 \rangle,$$

$$\langle 1, 0 \rangle \rightarrow_{189} \langle 0, 1 \rangle = \langle 0, 1 \rangle,$$

$$\langle 1, 0 \rangle \rightarrow_{189} \langle 1, 0 \rangle = \langle 0, 0 \rangle,$$

i.e., this operation, similarly to operation \rightarrow_{188} from [3], satisfies only a part of the basic properties of an implication. So, operation \rightarrow_{189} can be classified as a semi-implication.

Third, semi-implication \rightarrow_{189} generates the following negation

$$\neg^* \langle a, b \rangle = \langle a, b \rangle \rightarrow_{189} \langle 0, 1 \rangle = \langle 0, a \rangle,$$

i.e., the same negation as the one generated by implication \rightarrow_{188} from [3].

Therefore,

$$\neg^* \neg^* \langle a, b \rangle = \neg^* \langle 0, a \rangle = \langle 0, 0 \rangle. \quad (1)$$

Fourth, we see that

$$\begin{aligned} & (x \rightarrow_{189} y) \vee (y \rightarrow_{189} x) \\ &= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \vee (\langle c, d \rangle \rightarrow_{189} \langle a, b \rangle) \\ &= \langle bc, ad \rangle \vee \langle ad, bc \rangle \\ &= \langle ad, bc \rangle. \end{aligned}$$

Therefore, for this operation the above expression can not be an IFTP.

On the other hand, if we use the analogue of operation \vee in its new form \times_3 , we will obtain

$$\begin{aligned} & (x \rightarrow_{189} y) \times_3 (y \rightarrow_{189} x) \\ &= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \times_2 (\langle c, d \rangle \rightarrow_{189} \langle a, b \rangle) \\ &= \langle bc, ad \rangle \times_2 \langle ad, bc \rangle \\ &= \langle ad, b^2 c^2 \rangle, \end{aligned}$$

i.e., the situation is similar to above one.

Following [1], we check G. F. Rose's formula [6, 7], that has the form:

$$((\neg \neg x \rightarrow_{189} x) \rightarrow_{189} (\neg \neg x \vee \neg x)) \rightarrow_{189} (\neg \neg x \vee \neg x),$$

but compared to [3], here we use \neg^* instead of classical negation and will prove the following theorem.

Theorem 1. Rose's formula is an IFT.

Proof. Having in mind (1), we obtain sequentially:

$$\begin{aligned} & ((\neg^* \neg^* x \rightarrow_{189} x) \rightarrow_{189} (\neg^* \neg^* x \vee \neg^* x)) \rightarrow_{189} (\neg^* \neg^* x \vee \neg^* x) \\ &= ((\neg^* \neg^* \langle a, b \rangle \rightarrow_{189} \langle a, b \rangle) \rightarrow_{189} (\neg^* \neg^* \langle a, b \rangle \vee \neg^* \langle a, b \rangle)) \rightarrow_{189} (\neg^* \neg^* \langle a, b \rangle \vee \neg^* \langle a, b \rangle) \\ &= ((\langle 0, 0 \rangle \rightarrow_{189} \langle a, b \rangle) \rightarrow_{189} (\langle 0, 0 \rangle \vee \langle 0, a \rangle)) \rightarrow_{189} (\langle 0, 0 \rangle \vee \langle 0, a \rangle) \\ &= (\langle 0, 0 \rangle \rightarrow_{189} \langle 0, 0 \rangle) \rightarrow_{189} \langle 0, 0 \rangle \\ &= \langle 0, 0 \rangle \rightarrow_{189} \langle 0, 0 \rangle = \langle 0, 0 \rangle, \end{aligned}$$

which is an IFT. □

We obtain the same result, if we change operation \vee with operation \times_3 .

Fifth, following [1], we discuss the well-known Contraposition Law

$$(x \rightarrow_{189} y) \rightarrow_{189} (\neg y \rightarrow_{189} \neg x).$$

Theorem 2. The Contraposition Law is an IFT, but not a tautology as for classical negation, as well as for the new negation \neg^* .

Proof. Sequentially, we obtain:

$$\begin{aligned}
& (x \rightarrow_{189} y) \rightarrow_{189} (\neg y \rightarrow_{189} \neg x) \\
&= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \rightarrow_{189} (\neg \langle c, d \rangle \rightarrow_{189} \neg \langle a, b \rangle) \\
&= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \rightarrow_{189} (\langle d, c \rangle \rightarrow_{189} \langle b, a \rangle) \\
&= \langle bc, ad \rangle \rightarrow_{189} \langle bc, ad \rangle \\
&= \langle abcd, abcd \rangle,
\end{aligned}$$

which is an IFT.

$$\begin{aligned}
& (x \rightarrow_{189} y) \rightarrow_{189} (\neg^* y \rightarrow_{189} \neg^* x) \\
&= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \rightarrow_{189} (\neg^* \langle c, d \rangle \rightarrow_{189} \neg^* \langle a, b \rangle) \\
&= (\langle a, b \rangle \rightarrow_{189} \langle c, d \rangle) \rightarrow_{189} (\langle 0, c \rangle \rightarrow_{189} \langle 0, a \rangle) \\
&= \langle bc, ad \rangle \rightarrow_{189} \langle 0, 0 \rangle \\
&= \langle 0, 0 \rangle,
\end{aligned}$$

which is an IFT. □

Now, we check the validity of Klir and Yuan's axioms for fuzzy implications (marked by $I(x, y)$) [5], but in the intuitionistic fuzzy version from [1]:

Axiom A1 $(\forall x, y)(x \leq y \rightarrow (\forall z)(I(x, z) \geq I(y, z)))$,

Axiom A2 $(\forall x, y)(x \leq y \rightarrow (\forall z)(I(z, x) \leq I(z, y)))$,

Axiom A3 $(\forall y)(I(0, y) = 1)$,

Axiom A4 $(\forall y)(I(1, y) = y)$,

Axiom A5 $(\forall x)(I(x, x) = 1)$,

Axiom A6 $(\forall x, y, z)(I(x, I(y, z)) = I(y, I(x, z)))$,

Axiom A7 $(\forall x, y)(I(x, y) = 1 \text{ iff } x \leq y)$,

Axiom A8 $(\forall x, y)(I(x, y) = I(N(y), N(x)))$,

Axiom A9 I is a continuous function.

For our research, having in mind the specific forms of the intuitionistic fuzzy implication \rightarrow_{189} and following [1], we modify three of these axioms, as follows.

Axiom A3* $(\forall y)(I(0, y)$ is an IFT,

Axiom A5* $(\forall x)(I(x, x)$ is an IFT),

Axiom A7* $(\forall x, y)$ (if $x \leq y$, then, $I(x, y)$ is an IFT).

Theorem 5. Intuitionistic fuzzy implication \rightarrow_{189} satisfies axioms A1, A2, A3*, A5*, A6, A7*, A8 (for the classical negation N) and A9.

Proof. Let $x = \langle a, b \rangle, y = \langle c, d \rangle, z = \langle e, f \rangle$. We obtain sequentially. Let $x \leq y$. Then for $A1$ is valid:

$$I(x, z) = \langle be, af \rangle \geq \langle de, cf \rangle = I(y, z).$$

The checks for $A2$ is similar. For $A3^*$ we have

$$I(0, y) = \langle c, 0 \rangle,$$

i.e., $A3$ is not valid, but $A3^*$ is valid, while $A4$ is not valid, because

$$I(1, y) = \langle 0, d \rangle$$

is not an IFT. For $A5^*$ we obtain

$$I(x, x) = \langle a, b \rangle \rightarrow_{189} \langle a, b \rangle = \langle ab, ab \rangle.$$

Therefore, $A5^*$ is valid, while $A5$ is not. For $A6$ we have:

$$\begin{aligned} I(x, I(y, z)) &= \langle a, b \rangle \rightarrow_{189} (\langle c, d \rangle \rightarrow_{189} \langle e, f \rangle) \\ &= \langle a, b \rangle \rightarrow_{189} \langle de, cf \rangle \\ &= \langle bde, acf \rangle \\ &= \langle c, d \rangle \rightarrow_{189} \langle be, af \rangle \\ &= \langle c, d \rangle \rightarrow_{189} (\langle a, b \rangle \rightarrow_{189} \langle e, f \rangle) = I(y, I(x, z)), \end{aligned}$$

i.e., this axiom is valid.

Let $x \leq y$, i.e., $a \leq c$ and $b \geq d$. Then,

$$I(x, y) = \langle bc, ad \rangle,$$

which is an IFT. Therefore, $A7^*$ is valid, but $A7$ is not valid.

From

$$\begin{aligned} I(N(y), N(x)) &= \neg^* \langle c, d \rangle \rightarrow_{189} \neg^* \langle a, b \rangle \\ &= \langle 0, c \rangle \rightarrow_{189} \langle 0, a \rangle = \langle 0, 0 \rangle \neq I(x, y), \end{aligned}$$

it follows that $A8$ is not valid for the new negation, while,

$$\begin{aligned} I(N(y), N(x)) &= \neg \langle c, d \rangle \rightarrow_{189} \neg \langle a, b \rangle \\ &= \langle d, c \rangle \rightarrow_{189} \langle b, a \rangle = \langle bc, ad \rangle = I(x, y), \end{aligned}$$

i.e., $A8$ is valid.

Finally, obviously, $A9$ is valid. □

3 Conclusion

In a next step of this leg of research, other properties of the implication \rightarrow_{189} will be introduced and studied. For example, we will check the validity of axioms of intuitionistic logic, Kolmogorov's axioms and others.

References

- [1] Atanassov, K. (2017). *Intuitionistic Fuzzy Logics*. Springer, Cham.
- [2] Atanassov, K., Szmidt, E., & Kacprzyk, J. (2013). On intuitionistic fuzzy pairs, *Notes on Intuitionistic Fuzzy Sets*, 19(3), 1–13.
- [3] Atanassov, K., Szmidt, E., & Kacprzyk, J. (2017). Intuitionistic fuzzy implication \rightarrow_{189} . *Notes on Intuitionistic Fuzzy Sets*, 23(1), 6–13.
- [4] Atanassov, K., Szmidt, E., & Kacprzyk, J. (2017). Multiplicative type of operations over intuitionistic fuzzy pairs. *Proceedings of FQAS'17*, London, 21–22 June 2017, Springer (in press).
- [5] Klir, G., & Yuan, B. (1995). *Fuzzy Sets and Fuzzy Logic*. Prentice Hall, New Jersey.
- [6] Plisko, V. (2009). A survey of propositional realizability logic. *The Bulletin of Symbolic Logic*, 15(1), 1–42.
- [7] Rose, G. F. (1953). Propositional calculus and realizability. *Transactions of the American Mathematical Society*, 75, 1–19.