

# On the intuitionistic fuzzy implication $\rightarrow_{191}$

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*In memory of my friend Ivan Georgiev*

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**Abstract:** A new intuitionistic fuzzy implication,  $\rightarrow_{191}$ , is constructed. Some of its basic properties are studied.

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## 1 Introduction

In [3], 185 intuitionistic fuzzy implications are defined. Five other intuitionistic fuzzy implications are introduced in [4–10]. Now a new intuitionistic fuzzy implication will be given.

In some definitions we shall use functions  $sg$  and  $\overline{sg}$ :

$$sg(x) = \begin{cases} 1 & \text{if } x > 0 \\ 0 & \text{if } x \leq 0 \end{cases}, \quad \overline{sg}(x) = \begin{cases} 0 & \text{if } x > 0 \\ 1 & \text{if } x \leq 0 \end{cases}$$

Let everywhere intuitionistic fuzzy truth values of variables  $x$  and  $y$  be

$$x = \langle a, b \rangle, \quad y = \langle c, d \rangle.$$

In [3], for the variables  $x$  and  $y$  operation “conjunction” ( $\&$ ) is defined by:

$$V(x \& y) = \langle \min(a, c), \max(b, d) \rangle.$$

The pair  $\langle a, b \rangle$  is: a tautology if and only if (iff)  $a = 1$  and  $b = 0$ , and is an intuitionistic fuzzy tautology (IFT) iff  $a \geq b$ .

## 2 Main results

Let us define

$$x \rightarrow_{191} y = \langle a, b \rangle \rightarrow \langle c, d \rangle = \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b), \text{sg}(a - c) \text{sg}(d - b) \rangle.$$

For brevity, below we will write  $\rightarrow$  instead of  $\rightarrow_{191}$ .

First, we will show that the definition of the new implication is correct. Let  $a, b, c, d \in [0, 1]$  such that  $a + b \leq 1$  and  $c + d \leq 1$ . Then

$$0 \leq \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) \leq 1,$$

$$0 \leq \text{sg}(a - c) \text{sg}(d - b) \leq 1,$$

and if

$$X \equiv \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) + \text{sg}(a - c) \text{sg}(d - b),$$

then, we obtain sequentially. If  $a > c$ , then  $\text{sg}(a - c) = 1$  and  $\overline{\text{sg}}(a - c) = 0$ , i.e.,

$$X = 0 + \text{sg}(d - b) \leq 1.$$

If  $a \leq c$ , then,  $\text{sg}(a - c) = 0$  and  $\overline{\text{sg}}(a - c) = 1$ , i.e.,  $X = \overline{\text{sg}}(d - b) \leq 1$ .

Therefore, the definition of the new implication is correct. It generates the following negation.

$$\begin{aligned} \neg \langle a, b \rangle &= \langle a, b \rangle \rightarrow \langle 0, 1 \rangle = \langle \overline{\text{sg}}(a - 0) \overline{\text{sg}}(1 - b), \text{sg}(a - 0) \text{sg}(1 - b) \rangle \\ &= \langle \overline{\text{sg}}(a) \overline{\text{sg}}(1 - b), \text{sg}(a) \text{sg}(1 - b) \rangle = \langle \overline{\text{sg}}(1 - b), \text{sg}(a) \rangle \\ &= \begin{cases} \langle 1, 0 \rangle, & \text{if } a = 0 \text{ and } b = 1 \\ \langle 0, 0 \rangle, & \text{if } a = 0 \text{ and } b < 1 \\ & \text{or } a > 0 \text{ and } b = 1 \\ \langle 0, 1 \rangle, & \text{if } a > 0 \text{ and } b < 1 \end{cases} \end{aligned}$$

**Theorem 1.** *The new implication  $\rightarrow_{191}$ :*

- (a) *satisfies  $x \rightarrow x$  as a tautology;*
- (b) *satisfies  $x \rightarrow \neg \neg x$  as a tautology;*

(c) does not satisfy  $\neg\neg x \rightarrow x$  even as an IFT.

*Proof.* First, we see that

$$\begin{aligned}
\overline{\text{sg}}(1 - \text{sg}(a)) &= \begin{cases} \text{if } a = 0 : \overline{\text{sg}}(1) = 0 \\ \text{if } a > 0 : \overline{\text{sg}}(0) = 1 \end{cases} = \text{sg}(a), \\
\text{sg}(\overline{\text{sg}}(1 - b)) &= \begin{cases} \text{if } b = 1 : \text{sg}(1) = 1 \\ \text{if } b < 1 : \text{sg}(0) = 0 \end{cases} = \overline{\text{sg}}(1 - b), \\
\overline{\text{sg}}(a - \text{sg}(a)) &= \begin{cases} \text{if } a = 0 : \overline{\text{sg}}(0) = 1 \\ \text{if } a > 0 : \overline{\text{sg}}(0) = 1 \end{cases} = 1, \\
\text{sg}(a - \text{sg}(a)) &= \begin{cases} \text{if } a = 0 : \text{sg}(0) = 0 \\ \text{if } a > 0 : \text{sg}(a - 1) = 0 \end{cases} = 0, \\
\overline{\text{sg}}(\overline{\text{sg}}(1 - b) - b) &= \begin{cases} \text{if } b = 1 : \overline{\text{sg}}(0) = 1 \\ \text{if } b < 1 : \overline{\text{sg}}(0) = 1 \end{cases} = 1, \\
\overline{\text{sg}}(\text{sg}(a) - a) &= \begin{cases} \text{if } a = 0 : \overline{\text{sg}}(0) = 1 \\ \text{if } a = 0 : \overline{\text{sg}}(0) = 1 \\ \text{if } 0 < a < 1 : \overline{\text{sg}}(1) = 0 \end{cases}, \\
\overline{\text{sg}}(b - \overline{\text{sg}}(1 - b)) &= \begin{cases} \text{if } b = 0 : \overline{\text{sg}}(0) = 1 \\ \text{if } b = 1 : \overline{\text{sg}}(0) = 1 \\ \text{if } 0 < b < 1 : \overline{\text{sg}}(1) = 0 \end{cases},
\end{aligned}$$

Then

$$\neg\neg\langle a, b \rangle = \neg\langle \overline{\text{sg}}(1 - b), \text{sg}(a) \rangle = \langle \overline{\text{sg}}(1 - \text{sg}(a)), \text{sg}(\overline{\text{sg}}(1 - b)) \rangle$$

For (a) we obtain:

$$x \rightarrow x = \langle a, b \rangle \rightarrow \langle a, b \rangle = \langle \overline{\text{sg}}(0) \overline{\text{sg}}(0), \text{sg}(0) \text{sg}(0) \rangle = \langle 1, 0 \rangle.$$

For (b) we obtain:

$$\begin{aligned}
x \rightarrow \neg\neg x &= \langle a, b \rangle \rightarrow \neg\neg\langle a, b \rangle \\
&= \langle a, b \rangle \rightarrow \langle \overline{\text{sg}}(1 - \text{sg}(a)), \text{sg}(\overline{\text{sg}}(1 - b)) \rangle \\
&= \langle \overline{\text{sg}}(a - \text{sg}(a)) \overline{\text{sg}}(\overline{\text{sg}}(1 - b) - b), \text{sg}(a - \text{sg}(a)) \text{sg}(\overline{\text{sg}}(1 - b) - b) \rangle \\
&= \langle \overline{\text{sg}}(\overline{\text{sg}}(1 - b) - b), 0. \text{sg}(\overline{\text{sg}}(1 - b) - b) \rangle = \langle 1, 0 \rangle.
\end{aligned}$$

For (c) we see directly that

$$\begin{aligned}
\neg\neg x \rightarrow x &= \neg\neg\langle a, b \rangle \rightarrow \langle a, b \rangle \\
&= \langle \overline{\text{sg}}(1 - \text{sg}(a)), \text{sg}(\overline{\text{sg}}(1 - b)) \rangle \rightarrow \langle a, b \rangle \\
&= \langle \overline{\text{sg}}(\overline{\text{sg}}(1 - \text{sg}(a)) - a) \overline{\text{sg}}(b - \text{sg}(\overline{\text{sg}}(1 - b))), \text{sg}(\overline{\text{sg}}(1 - \text{sg}(a)) - a) \text{sg}(b - \text{sg}(\overline{\text{sg}}(1 - b))) \rangle \\
&= \langle \overline{\text{sg}}(\overline{\text{sg}}(1 - \text{sg}(a)) - a) \overline{\text{sg}}(b - \text{sg}(\overline{\text{sg}}(1 - b))), \text{sg}(\overline{\text{sg}}(1 - \text{sg}(a)) - a) \text{sg}(b - \text{sg}(\overline{\text{sg}}(1 - b))) \rangle
\end{aligned}$$

(for  $0 < a, b < 1$ )

$$= \langle 0, 1 \rangle,$$

i.e., it is not an IFT. □

**Theorem 2.** *The new implication  $\rightarrow_{191}$ :*

- (a) *satisfies Modus Ponens in tautological sense,*
- (b) *does not satisfy Modus Ponens in the IFT-sense,*
- (c) *satisfies for every two variables  $x$  and  $y$ ,*

$$(x \& (x \rightarrow y)) \rightarrow y$$

*in the IFT-sense.*

*Proof.* (a) Let  $\langle a, b \rangle$  be a tautology, i.e.,  $a = 1$  and  $b = 0$  and let  $\langle a, b \rangle \rightarrow \langle c, d \rangle$  be a tautology, i.e.,  $\langle \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b) = 1$  and  $\text{sg}(a-c) \text{sg}(d-b) = 0$ . Then  $1 = \overline{\text{sg}}(1-c) \overline{\text{sg}}(d-0) = \overline{\text{sg}}(1-c) \overline{\text{sg}}(d)$ , i.e.  $1 = \overline{\text{sg}}(1-c)$  and hence  $c = 1$  and  $d = 0$ . Therefore  $\langle c, d \rangle$  is a tautology.

(b) Let  $\langle a, b \rangle = \langle 0, 0 \rangle$ , i.e., an IFT. Then,  $\langle \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b), \text{sg}(a-c) \text{sg}(d-b) \rangle$  will be an IFT, e.g., for  $\langle c, d \rangle = \langle 0.1, 0.2 \rangle$ , but the last pair is not an IFT.

(c) we have sequentially:

$$\begin{aligned} & V((x \& (x \rightarrow y)) \rightarrow y) \\ &= (\langle a, b \rangle \& \langle \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b), \text{sg}(a-c) \text{sg}(d-b) \rangle) \rightarrow \langle c, d \rangle \\ &= \langle \min(a, \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b)), \max(b, \text{sg}(a-c) \text{sg}(d-b)) \rangle \rightarrow \langle c, d \rangle \\ &= \langle \overline{\text{sg}}(\min(a, \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b)) - c) \overline{\text{sg}}(d - \max(b, \text{sg}(a-c) \text{sg}(d-b))), \\ & \quad \text{sg}(\min(a, \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b)) - c) \text{sg}(d - \max(b, \text{sg}(a-c) \text{sg}(d-b))) \rangle. \end{aligned}$$

Let

$$\begin{aligned} X &\equiv \overline{\text{sg}}(\min(a, \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b)) - c) \overline{\text{sg}}(d - \max(b, \text{sg}(a-c) \text{sg}(d-b))) \\ &\quad - \text{sg}(\min(a, \overline{\text{sg}}(a-c) \overline{\text{sg}}(d-b)) - c) \text{sg}(d - \max(b, \text{sg}(a-c) \text{sg}(d-b))). \end{aligned}$$

If  $a > c$ , then

$$\begin{aligned} X &= \overline{\text{sg}}(\min(a, 0) - c) \overline{\text{sg}}(d - \max(b, \text{sg}(d-b))) - \text{sg}(\min(a, 0) - c) \text{sg}(d - \max(b, \text{sg}(d-b))) \\ &= \overline{\text{sg}}(-c) \overline{\text{sg}}(d - \max(b, \text{sg}(d-b))) - \text{sg}(-c) \text{sg}(d - \max(b, \text{sg}(d-b))) \\ &= \overline{\text{sg}}(d - \max(b, \text{sg}(d-b))) \geq 0. \end{aligned}$$

If  $a \leq c$ , then

$$\begin{aligned} X &= \overline{\text{sg}}(\min(a, \overline{\text{sg}}(d-b)) - c) \overline{\text{sg}}(d - \max(b, 0)) - \text{sg}(\min(a, \overline{\text{sg}}(d-b)) - c) \text{sg}(d - \max(b, 0)) \\ &= \overline{\text{sg}}(\min(a, \overline{\text{sg}}(d-b)) - c) \overline{\text{sg}}(d-b) - \text{sg}(\min(a, \overline{\text{sg}}(d-b)) - c) \text{sg}(d-b). \end{aligned}$$

If  $d > b$ , then

$$X = \overline{\text{sg}}(\min(a, 0) - c).0 - \text{sg}(\min(a, 1) - c) = 0 - \text{sg}(a - c) = 0.$$

If  $d \leq b$ , then

$$X = \overline{\text{sg}}(\min(a, 1) - c) - \text{sg}(\min(a, 1) - c).0 = \overline{\text{sg}}(a - c) = 1,$$

i.e., the expression is an IFT. □

Following [3], let us define for  $x$

$$\Box x = \langle a, 1 - a \rangle.$$

**Theorem 3.** *The new implication  $\rightarrow_{191}$  satisfies the formula*

$$\Box(x \rightarrow y) \rightarrow (\Box x \rightarrow \Box y) \quad (*)$$

as a tautology.

*Proof.* We obtain sequentially

$$\begin{aligned} & \Box(x \rightarrow y) \rightarrow (\Box x \rightarrow \Box y) \\ &= \Box(\langle a, b \rangle \rightarrow \langle c, d \rangle) \rightarrow (\Box \langle a, b \rangle \rightarrow \Box \langle c, d \rangle) \\ &= \Box(\langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b), \text{sg}(a - c) \text{sg}(d - b) \rangle) \rightarrow (\langle a, 1 - a \rangle \rightarrow \langle c, 1 - c \rangle) \\ &= \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b), 1 - \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) \rangle \rightarrow \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(1 - c - 1 + a), \text{sg}(a - c) \text{sg}(1 - c - 1 + a) \rangle \\ &= \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b), 1 - \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) \rangle \rightarrow \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(a - c), \text{sg}(a - c) \text{sg}(a - c) \rangle \\ & \text{(because } \overline{\text{sg}}(p) \overline{\text{sg}}(p) = \overline{\text{sg}}(p) \text{ and } \text{sg}(p) \text{sg}(p) = \text{sg}(p) \text{ for each } p \in [0, 1]) \\ &= \langle \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b), 1 - \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) \rangle \rightarrow \langle \overline{\text{sg}}(a - c), \text{sg}(a - c) \rangle \\ &= \langle \overline{\text{sg}}(\overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) - \overline{\text{sg}}(a - c)) \overline{\text{sg}}(\text{sg}(a - c) - 1 + \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b)), \\ & \quad \text{sg}(\overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) - \overline{\text{sg}}(a - c)) \text{sg}(\text{sg}(a - c) - 1 + \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b)) \rangle. \end{aligned}$$

Let

$$\begin{aligned} X &= \langle \overline{\text{sg}}(\overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) - \overline{\text{sg}}(a - c)) \overline{\text{sg}}(\text{sg}(a - c) - 1 + \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b)), \\ & \quad \text{sg}(\overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b) - \overline{\text{sg}}(a - c)) \text{sg}(\text{sg}(a - c) - 1 + \overline{\text{sg}}(a - c) \overline{\text{sg}}(d - b)) \rangle. \end{aligned}$$

If  $a > c$ , then

$$\begin{aligned} x &= \langle \overline{\text{sg}}(0. \overline{\text{sg}}(d - b) - 0) \overline{\text{sg}}(1 - 1 + 0. \overline{\text{sg}}(d - b)), \text{sg}(0. \overline{\text{sg}}(d - b) - 0) \text{sg}(1 - 1 + \overline{\text{sg}}(d - b)) \rangle \\ &= \langle \overline{\text{sg}}(0) \overline{\text{sg}}(0), \text{sg}(0) \text{sg}(\overline{\text{sg}}(d - b)) \rangle = \langle 1, 0 \rangle. \end{aligned}$$

If  $a \leq c$ , then

$$X = \langle \overline{\text{sg}}(\overline{\text{sg}}(d - b) - 1) \overline{\text{sg}}(-1 + \overline{\text{sg}}(d - b)), \text{sg}(\overline{\text{sg}}(d - b) - 1) \text{sg}(-1 + \overline{\text{sg}}(d - b)) \rangle$$

(because  $\overline{\text{sg}}(d - b) \leq 1$ )

$$X = \langle \overline{\text{sg}}(0), \text{sg}(\overline{\text{sg}}(d - b) - 1) \rangle = \langle 1, 0 \rangle.$$

Therefore formula  $(*)$  is a tautology. Hence, it is an IFT, too. □

### 3 Conclusion

In this paper we have introduced the 191-st intuitionistic fuzzy implication, and have shown that it has intuitionistic but not classical behaviour. Several of its properties were formulated and proved. In a future research, new properties of the new implications will be studied.

**Open Problem:** In [2] six Cartesian products are defined (the last of them was published in [1]). For which index  $j$  the  $j$ -th Cartesian product satisfies equality

$$(A \rightarrow_{191} B) \times_j C = (A \times_j C) \rightarrow_{191} (B \times_j C)$$

for every three IFSs  $A, B, C$ ?

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